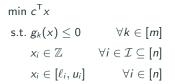
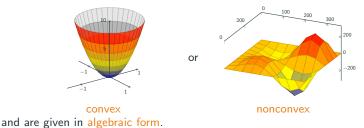
Nonlinear constraints in SCIP

Ksenia Bestuzheva and Stefan Vigerske SCIP Online Workshop 2020 · June 3, 2020

Mixed-Integer Nonlinear Programming



• The functions $g_k : [\ell, u] \to \mathbb{R}$ can be



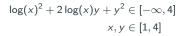
• SCIP solves MINLPs by spatial Branch & Bound.

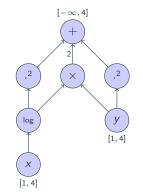
The "classical" framework for (MI)NLP in SCIP

Expression trees and graphs

cons_nonlinear stores algebraic structure of nonlinear constraints in one directed acyclic graph:

- nodes: variables, operations, constraints
- arcs: flow of computation





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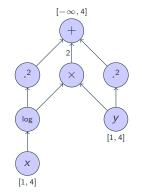
- nodes: variables, operations, constraints
- arcs: flow of computation

Operators:

- variable index, constant
- +, -, *, ÷
- \cdot^2 , $\sqrt{\cdot}$, \cdot^p $(p \in \mathbb{R})$, \cdot^n $(n \in \mathbb{Z})$, $x \mapsto x |x|^{p-1}$ (p > 1)
- exp, log
- min, max, abs
- ∑, ∏, affine-linear, quadratic, signomial
- (user)

$$\log(x)^{2} + 2\log(x)y + y^{2} \in [-\infty, 4]$$

x, y \in [1, 4]



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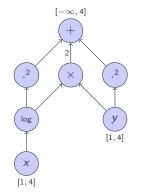
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Additional constraint handler: quadratic, abspower $(x \mapsto x|x|^{p-1}, p > 1)$, SOC

 $log(x)^{2} + 2 log(x)y + y^{2} \in [-\infty, 4]$ $x, y \in [1, 4]$

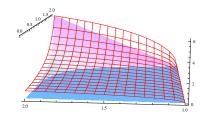


Goal: Reformulate constraints such that only elementary cases (convex, concave, odd power, quadratic) remain.

Example:

$$g(x) = \sqrt{\exp(x_1^2) \ln(x_2)}$$

 $x_1 \in [0, 2], \quad x_2 \in [1, 2]$



- reformulates constraints by introducing new variables and new constraints
- other constraint handler can participate

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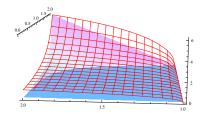
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Reformulation:

 $g = \sqrt{y_1}$ $y_1 = y_2 y_3$ $y_2 = \exp(y_4)$ $y_3 = \ln(x_2)$ $y_4 = x_1^2$



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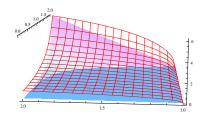
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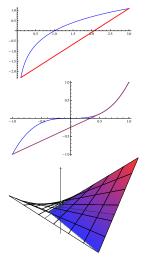
 $g = \sqrt{y_1}$ $y_1 = y_2 y_3 \qquad [0, \ln(2)e^4]$ $y_2 = \exp(y_4) \qquad [1, e^4]$ $y_3 = \ln(x_2) \qquad [0, \ln(2)]$ $y_4 = x_1^2 \qquad [0, 4]$



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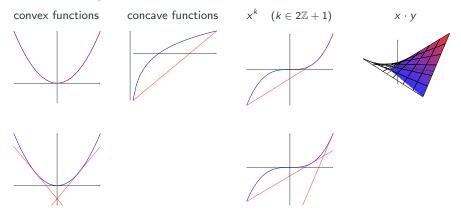
Bounding: LP relaxation

- relaxing integrality
- convexifying non-convexities
- linearizing nonlinear convexities



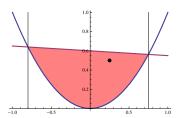
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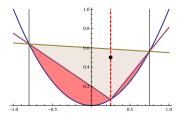
- relaxing integrality
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- 1. fractional integer variables
- 2. variables in violated nonconvex constraints, because variable bounds determine the convex relaxation, e.g.,

$$x^2 \leq \ell^2 + rac{u^2 - \ell^2}{u - \ell} (x - \ell) \quad \forall x \in [\ell, u].$$

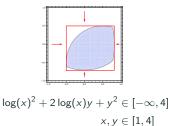


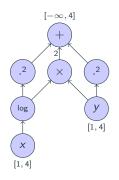


 Many constraint handler in SCIP implement a fairly cheap bound tightening (aka. domain propagation) method to infer tighter variable bounds from constraints and current variable bounds.

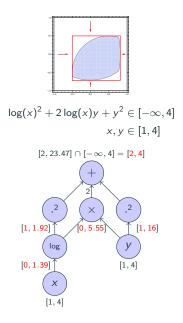


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- cons_nonlinear utilizes the expression graph and interval arithmetic

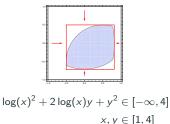


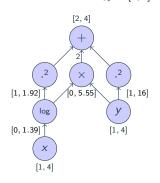


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 - compute bounds on intermediate nodes (top-down)

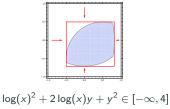


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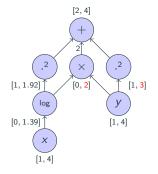




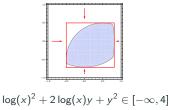
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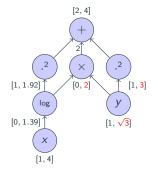
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Primal Heuristics:

- NLP solving: subnlp, nlpdiving, multistart, mpec
- MINLP solving: LNS heuristics (RENS, RINS, DINS, etc.)
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Tighter Relaxations:

- second-order cone upgrade of quadratic constraints
- adding KKT reformulation (using SOS1) for QPs
- $x \cdot y$ over 2D projections of the LP relaxation
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More Bound Tightening:

- Optimization-Based Bound Tightening: min / max x_i w.r.t. LP relaxation
- nl. Optimization-Based Bound Tightening: min / max x_i w.r.t. convex NLP relaxation (off by default)

A MINLP can be input via

- File readers: FlatZinc*, LP*, MPS*, OSiL, PIP[†], ZIMPL
- Interfaces: AMPL, C, GAMS, Java*, Julia/JuMP, Matlab (via OPTI Toolbox), Python

SCIP can utilize this software for MINLP solving:

- NLP Solvers: Ipopt, FilterSQP, WORHP
- Automatic Differentiation: CppAD

*quadratic only
t polynomial only

Example: Circle Packing

A new framework for NLP in SCIP (work in progress) by K. Bestuzheva, B. Müller, F. Serrano, S. Vigerske, F. Wegscheider

Problem with current implementation

Consider

min z s.t. $\exp(\ln(1000) + 1 + xy) \le z$ $x^2 + y^2 \le 2$ x = -1 y = 1z = 1000

An optimal solution:

Problem with current implementation

Consider	min z s.t. exp(ln(1000) + 1 - $x^2 + y^2 \le 2$	$+xy) \leq z$	An optimal solution: x = -1 y = 1 z = 1000	
SCIP reports				
<pre>SCIP Status : problem is solved [optimal solution found] Solving Time (sec) : 0.08 Solving Nodes : 5 Primal Bound : +9.99999656552062e+02 (3 solutions) Dual Bound : +9.99999656552062e+02 Gap : 0.00 % [nonlinear] <e1>: exp((7.9077552789821368151 +1 (<x> * <y>)))-1<z>[C] <= 0; violation: right hand side is violated by 0.000673453314561812 best solution is not feasible in original problem</z></y></x></e1></pre>				
x	-	1.00057454873626	(obj:0)	

У	0.999425451364613	(obj:0)
z	999.999656552061	(obj:1)

Reformulation takes apart exp(ln(1000) + 1 + xy), thus SCIP actually solves

min z

s.t.
$$\exp(w) \le z$$

 $\ln(1000) + 1 + xy = w$
 $x^2 + y^2 \le 2$

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 $\begin{array}{ll} \min z & \mbox{Violation} \\ \mbox{s.t. } \exp(w) \leq z & \mbox{0.4659} \cdot 10^{-6} \leq \mbox{numerics/feastol} \, \checkmark \\ & \mbox{ln}(1000) + 1 + x \, y = w & \mbox{0.6731} \cdot 10^{-6} \leq \mbox{numerics/feastol} \, \checkmark \\ & \mbox{x}^2 + y^2 \leq 2 & \mbox{0.6602} \cdot 10^{-6} \leq \mbox{numerics/feastol} \, \checkmark \\ \end{array}$

Solution (found by <relaxation>):

x = -1.000574549

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- \Rightarrow Explicit reformulation of constraints ...
 - ... loses the connection to the original problem.
 - ... loses distinction between original and auxiliary variables. Thus, we may branch on auxiliary variables.
 - ... prevents simultaneous exploitation of overlapping structures.

Everything is an expression.

- ONE constraint handler: cons_expr
- represent all nonlinear constraints in one expression graph (DAG)

 $lhs \leq expression-node \leq rhs$

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Do not reformulate constraints.

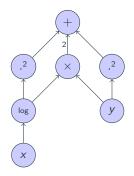
• introduce auxiliary variables for the relaxation only

Constraint:

$$\log(x)^2 + 2\log(x)y + y^2 \le 4$$

This formulation is used to

- check feasibility,
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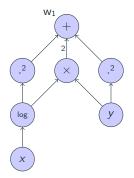
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(Implicit) Reformulation:

$$w_1 \le 4$$

 $\log(x)^2 + 2\log(x)y + y^2 = w_1$



Enforcement

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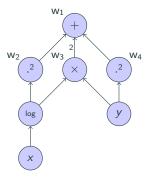
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(Implicit) Reformulation:

$$w_1 \le 4$$

 $w_2 + 2w_3 + w_4 = w_1$
 $\log(x)^2 = w_2$
 $\log(x)y = w_3$
 $y^2 = w_4$



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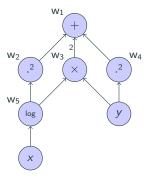
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 $w_5y = w_3$
 $y^2 = w_4$
 $\log(x) = w_5$



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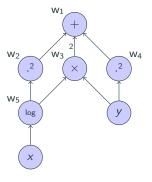
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 $\log(x) = w_5$



Used to construct LP relaxation.

Each operator type $(+, \times, pow, etc.)$ is implemented by an **expression** handler, which can provide a number of callbacks:

- evaluate and differentiate expression w.r.t. operands
- interval evaluation and tighten bounds on operands
- provide linear under- and over-estimators
- inform about curvature, monotonicity, integrality
- simplify, compare, print, parse, hash, copy, etc.

Expression handler are like other SCIP plugins, thus new ones can be added by users.

```
min z s.t. \exp(\ln(1000) + 1 + xy) \le z, x^2 + y^2 \le 2
```

Classic:

```
presolving (5 rounds: 5 fast, 1 medium, 1 exhaustive):
  0 deleted vars, 0 deleted constraints, 1 added constraints,...
  0 implications, 0 cliques
  presolved problem has 4 variables (0 bin, 0 int, 0 impl, 4 cont)
  and 3 constraints
        2 constraints of type <quadratic>
        1 constraints of type <quadratic>
```

[...]

SCIP Status	:	problem is solved [optimal solution found]
Solving Time (sec)	:	0.08
Solving Nodes	:	5
Primal Bound	:	+9.99999656552062e+02 (3 solutions)
Dual Bound	:	+9.99999656552062e+02
Gap	:	0.00 %

 $\label{eq:continuous} \begin{array}{l} (e1>: exp((7.90776 + (<x> * < y>))))-1<z>[C] <= 0; \\ \mbox{violation: right hand side is violated by } 0.000673453314561812 \\ \mbox{best solution is not feasible in original problem} \end{array}$

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Classic:

New:

presolving (5 rounds: 5 fast, 1 medium, 1 exhaustive): presolving (3 rounds: 3 fast, 1 medium, 1 exhaustive): 0 deleted vars, 0 deleted constraints, 1 added constraints, 0.deleted vars, 0 deleted constraints, 0 added constraints,... 0 implications, 0 cliques presolved problem has 4 variables (0 bin, 0 int, 0 impl, 4 prombleved problem has 3 variables (0 bin, 0 int, 0 impl, 3 cont) and 3 constraints of type <quadratic> 1 constraints of type <quadratic> 2 constraints of type <nolinear>

```
[...]
```

[...]

SCIP Status	: problem is solved [optimal solution for	Status	: problem is solved [optimal solution found]						
Solving Time (sec)	: 0.08	Solving Time (sec)	: 0.47						
Solving Nodes	: 5	Solving Nodes	: 15						
Primal Bound	: +9.99999656552062e+02 (3 solutions)	Primal Bound	: +9.999999949950021e+02 (2 solutions)						
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Constraint:
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Smarter reformulation:

• Recognize that $\log(x)^2 + 2\log(x)y + y^2$ is convex in $(\log(x), y)$.

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Smarter reformulation:

• Recognize that $\log(x)^2 + 2\log(x)y + y^2$ is convex in $(\log(x), y)$.

 \Rightarrow Introduce auxiliary variable for log(x) only.

$$w^2 + 2wy + y^2 \le 4$$
$$\log(x) = w$$

Handle $w^2 + 2wy + y^2 \le 4$ as convex constraint ("gradient-cuts").

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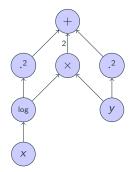
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Nonlinearity Handler (nlhdlrs):

- Adds additional separation and propagation algorithms for structures that can be identified in the expression graph.
- Attached to nodes in expression graph, but does not *define* expressions or constraints.
- Examples: quadratics, convex subexpressions, vertex-polyhedral

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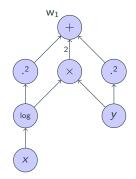


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 $w_1 \leq 4$

1. Add auxiliary variable w_1 for root.

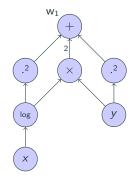


- Nodes in the expression graph can have one or several nlhdlrs attached.
- At beginning of solve, detection callbacks are run **only** for nodes that have auxiliary variable. Detection callback may add auxiliary variables.

Constraint: $\log(x)^2 + 2\log(x)y + y^2 \le 4$

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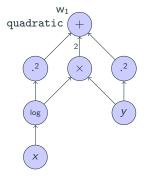


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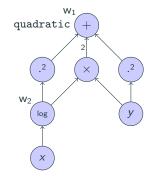
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$$w_1 \leq 4 \\ w_2^2 + 2w_2y + y^2 \leq w_1 \quad \texttt{[nlhdlr_quadratic]}$$

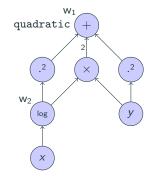
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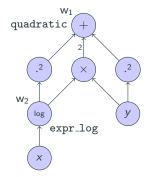


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$$w_1 \leq 4$$

 $w_2^2 + 2w_2y + y^2 \leq w_1 \quad [nlhdlr_quadratic]$
 $\log(x) = w_2 \quad [expr_log]$

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 - nlhdlr_quadratic detects a convex quadratic structure and signals success.
 - nlhdlr_quadratic adds an auxiliary variable w₂ for log node.
- 3. Run detect of all nlhdlrs on log node.
 - No specialized nlhdlr signals success. The expression handler will be used.



Current status

Available features:

- Handler for quadratic subexpressions
- Handler for second-order cone structures
- Handler for convex and concave subexpressions
- Handler for functions on semi-continuous variables (perspective formulations)
- Handler for bilinear terms $(x \cdot y \text{ over 2D projection of LP relaxation})$
- RLT (Reformulation-Linearization Technique) separator for bilinear terms
- Separator for SDP cuts on 2×2 principal minors of $X xx^{\mathsf{T}} \succeq 0$
- Linearization of products of binary variables
- Symmetry detection
- Support for operators cos, sin, entropy
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Main open tasks before release:

- check performance "outliers"
- replacing remaining classic code by new one (in particular NLP relaxation and NLP solver interfaces)
- documentation "?"

Nonlinear constraints in SCIP

Ksenia Bestuzheva and Stefan Vigerske SCIP Online Workshop 2020 · June 3, 2020